# Intrinsic Triangulations in Geometry Processing 

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## Intrinsic triangles




## Triangle meshes can be very frustrating



## Triangle meshes can be very frustrating



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## Outline

I. Preliminaries


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I. Preliminaries
II. Data structures for intrinsic triangulations

[ G., Sharp, \& Crane. 2021.
Integer coordinates for intrinsic geometry processing. ACM TOG ]

- Integer data structure for intrinsic triangulations
- Quality guarantees for intrinsic remeshing


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I. Preliminaries
II. Data structures for intrinsic triangulations

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- Integer data structure for intrinsic triangulations
- Quality guarantees for intrinsic remeshing
III. Simplification of intrinsic triangulations

[ Liu, G., Chislett, Sharp, Jacobson \& Crane. 2023. Surface Simplification using Intrinsic Error Metrics. ACM TOG ]
- First algorithm for intrinsic simplification
- New distortion measurement via intrinsic curvature error


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## III. Simplification of intrinsic triangulations


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- First algorithm for intrinsic simplification
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## IV. Discrete conformal

 maps \& uniformization
[ G., Springborn, \& Crane. 2021. Discrete conformal equivalence of polyhedral surfaces. ACM TOG ]

- Data structure for ideal hyperbolic polyhedra
- Interpolation for hyperbolic isometries
- Careful treatment of numerics


## I. Preliminaries



## Triangulations

## Definition

A (surface) triangulation is a manifold 2-dimensional cell complex $T=(V, E, F)$ whose faces are all triangles

- May be irregular (e.g., two edges of a face may be glued together)



## Intrinsic and extrinsic triangulations

## Definition

An extrinsic triangulation is a triangulation equipped with a piecewise-linear embedding into $\mathbb{R}^{3}$, i.e., vertex positions $p: V \rightarrow \mathbb{R}^{3}$

## Definition

An intrinsic triangulation is a triangulation equipped with positive edge lengths $\ell: E \rightarrow \mathbb{R}_{>0}$ satisfying the triangle inequality within each face


## Correspondence



- Common case: intrinsic triangulation on top of extrinsic triangulation
- i.e. isometric or at least homeomorphic to extrinsic triangulation
- The correspondence is the homeomorphism mapping between them



## Common subdivision



## Definition

The common subdivision of two triangulations $T_{1}, T_{2}$ is the coarsest polygonal complex $\mathcal{S}$ such that all faces of $T_{1}$ or $T_{2}$ are unions of faces of $\mathcal{S}$

Intuitively, the result of cutting $T_{1}$ along edges of $T_{2}$


## The space of intrinsic triangulations is large



## Delaunay triangulations

I. Preliminaries

- Planar Delaunay triangulations have many nice properties:
- Essentially unique, maximize angles lexicographically, minimize spectrum lexicographically, smoothest interpolation, positive cotan weights...
- Characterized by empty circumcircle condition


$$
\alpha+\beta \leq \gamma+\delta
$$



## Intrinsic Delaunay triangulations


I. Preliminaries

- [Indermitte, Liebling, Troyanov \& Clemençon 2001, Bobenko \& Springborn 2007]: empty intrinsic circumcircles
- Maintain most nice properties. [Sharp, G. \& Crane 2021; §4.1.1]
- Compute by a simple algorithm:
- Flip any non-Delaunay edge until none remain



## Intrinsic Delaunay triangulations provide good function spaces


I. Preliminaries


## A brief history of intrinsic triangulations


I. Preliminaries


Foundations: [Alexandrov 1948; Regge 1961]
Geometry Processing: [Fisher, Springborn, Bobenko \& Schröder 2006; Bobenko \& Springborn 2007, Bobenko \& Izmestiev 2008; Sun, Wu, Gu \& Luo 2015; Sharp, Soliman \& Crane 2019; Fumero, Möller \& Rodolà 2020; Gillespie, Springborn \& Crane 2021; Finnendahl, Schwartz \& Alexa 2023]

## II. Integer Coordinates for Intrinsic Triangulations


$\left[\begin{array}{c}\text { G., Sharp, \& Crane. 2021. Integer coordinates for intrinsic geometry } \\ \text { processing. ACM Transactions on Graphics }\end{array}\right]$

## Correspondence data structures

Overlay Mesh

[Fisher, Springborn, Bobenko \& Schröder 2006]

- Explicit mesh of common subdivision
- Edge flips nonlocal \& expensive
- No further operations


## Signposts

Integer coordinates combine the best of both worlds
[Sharp, Soliman \& Crane 2019]

- Floating point signpost vectors at vertices
- Supports many local mesh operations
- Common subdivision connectivity may be invalid


## The integer coordinates data structure

II. Integer coordinates for intrinsic triangulations


[^0]
## Normal coordinates


II. Integer coordinates for intrinsic triangulations

Foundations: [Kneser 1929; Haken 1961]
Computational Topology: [Schaefer+ 2008; Erickson \& Nayyeri 2013]

## Normal coordinates


II. Integer coordinates for intrinsic triangulations

## Slight complication

- Standard setting: homotopy classes of closed curves on a topological surface (or closed surfaces in a topological 3-manifold)
- Our setting: edges of a geodesic triangulation on a Riemannian manifold

Foundations: [Kneser 1929; Haken 1961]
Computational Topology: [Schaefer+ 2008; Erickson \& Nayyeri 2013]

## Encoding a curve with normal coordinates


II. Integer coordinates for intrinsic triangulations

- Just count intersections


## Rules

1. No self-crossings
2. No U-turns
(also curves may only start or end
 at vertices of the triangulation) automatically satisfied for geodesic triangulations


## How much do normal coordinates tell us?

II. Integer coordinates for intrinsic triangulations

- Represents curve up to homotopy (on the surface punctured at vertices)
- Equivalently, encodes a sequence of triangles



## Reconstructing the curve



## Rules

1. No self-crossings
2. No U-turns


## Reconstructing the curve

- Normality conditions determine curves within each triangle



## Finding the exact curve geometry

- So far: triangle strip
- True curve is geodesic
- Lay out in plane to find exact curve
- Normal coordinates determine edges exactly



## Collections of Curves

- e.g. edges of a triangulation
- Could store multiple sets of normal coordinates
- Expensive
- Instead, just store one set of normal coordinates

Store entire triangulation using one integer per edge

## The integer coordinates data structure

II. Integer coordinates for intrinsic triangulations

normal coordinates


## Normal coordinates are not enough for correspondence

II. Integer coordinates for intrinsic triangulations

- How do you tell what edge you have traced?



## Normal coordinates are not enough for correspondence

II. Integer coordinates for intrinsic triangulations

- How do you tell what edge you have traced?



## Normal coordinates are not enough for correspondence

II. Integer coordinates for intrinsic triangulations

- How do you tell what edge you have traced?
- Disambiguate with roundabouts



## Roundabouts


II. Integer coordinates for intrinsic triangulations

- Each black edge stores a pointer to the next yellow curve
- Resolves all ambiguity



## Data structure operations



II. Integer coordinates for intrinsic triangulations - connectivity changes

- Supports a variety of connectivity changes:

edge flips

vertex insertion

flat vertex removal


## Edge flips

II. Integer coordinates for intrinsic triangulations

Normal Coordinates


$$
n_{k l}=\max \left(n_{k i}+n_{l j}, n_{j k}+n_{l i}\right)-n_{i j}
$$

$$
r_{k l}=r_{k i}+n_{k i}^{-}+\max \left(0, n_{i l}^{+}-n_{l k}^{+}-n_{k i}^{+}\right)
$$

(if there are no endpoints)

## Edge flips

II. Integer coordinates for intrinsic triangulations

## Normal Coordinates



$$
\begin{aligned}
& \left.n_{k l}=c_{l}^{j k}+c_{k}^{i j}+\frac{1}{2}\left|c_{j}^{i l}-c_{j}^{k i}\right|+\frac{1}{2} \right\rvert\, c_{i}^{l j}-c_{i}^{j k} \\
& -\frac{1}{2} e_{l}^{j i}-\frac{1}{2} e_{k}^{i j}+e_{i}^{l j}+e_{i}^{j k}+e_{j}^{i l}+e_{j}^{k i}+n_{i j}^{-} \\
& \text {( general case ) } \\
& 2 c_{k}^{i j}:=\underset{-\rho_{j}^{j k}-e^{k i}}{\max \left(0, n_{j i}^{+}+n_{k i}^{+}-n_{i j}^{+}\right)} \\
& e_{k}^{i j}:=\max \left(0, n_{i j}^{+}-n_{j k}^{+}-n_{k i}^{+}\right)
\end{aligned}
$$

## Key takeaway:

closed form flip formulas, independent of geometry

## Roundabouts



## Vertex insertion

_- edge $\rightleftharpoons$ curve

- Unlike classic normal coordinates, depends on geometry
- Not a computational challenge:

1. Locate curves via normal coordinates
2. Count intersections
3. Update roundabouts


## Vertex removal

## _- edge

 $\rightleftharpoons$ curveII. Integer coordinates for intrinsic triangulations

- Only remove inserted vertices
- Strategy: reduce to degree-3 case



## Vertex removal

II. Integer coordinates for intrinsic triangulations

- Only remove inserted vertices
- Strategy: reduce to degree-3 case
II. Integer coordinates for intrinsic triangulations


## Applications

## Delaunay refinement for planar meshing

II. Integer coordinates for intrinsic triangulations

- Crucial tool in 2D - remesh with guaranteed quality bounds [Chew 1993; Shewchuk 1997]



## Intrinsic Delaunay refinement

- Intrinsic retriangulation algorithm proposed by [Sharp, Soliman \& Crane 2019]


## Theorem [G., Sharp \& Crane 2021]

Let $M$ be a mesh without boundary whose cone angles are all at least $60^{\circ}$. Then intrinsic Delaunay refinement produces a Delaunay mesh with triangle corner angles at least $30^{\circ}$

II. Integer coordinates for


## Common subdivisions of intrinsic Delaunay refinements


II. Integer coordinates for intrinsic triangulations

- Integer coordinates can be crucial to recovering the common subdivision



## Intrinsic Delaunay refinement of meshes with boundary

Integer coordinates for intrinsic triangulations

- Extend algorithm to meshes with boundary


ThingilD 48352

## Intrinsic Delaunay refinement - validation

II. Integer coordinates for intrinsic triangulations

- Compute refinements \& common subdivisions for Thingi10k dataset [Zhou \& Jacobson 2016]
- 7696 manifold meshes
- < 1s on most meshes; only took > 1m on 6 meshes
- $100 \%$ success rate for refinement $\&$ common subdivision
- [Sharp, Soliman \& Crane 2019] succeed on only $69.1 \%$ of meshes



## Application: PDE-Based Geometry Processing



II. Integer coordinates for intrinsic triangulations
mean error: $2 \%$
result on Delaunay refinement

## Application: Flip-Based Geodesic Paths

- FlipOut [Sharp \& Crane 2020]:
- computes geodesic paths via edge flips



## Try it out yourself


https://github.com/MarkGillespie/intrinsic-triangulations-demo

## III. Simplifying Intrinsic Triangulations


[Liu, G., Chislett, Sharp, Jacobson, \& Crane. 2023. Surface Simplification using Intrinsic Error Metrics. ACM Transactions on Graphics

## Exact geometry preservation: a blessing and a curse


mean error: $2 \%$
III. Intrinsic simplification

- motivation




## Coarse meshes can be perfectly adequate

 - motivation
## Coarse meshes can be perfectly adequate

III. Intrinsic simplification

23.14 s


## Traditional goal: extrinsic simplification

III. Intrinsic simplification - motivation

- Find a coarse mesh close in space to the original
- Often designed to optimize for visual fidelity



## Intrinsic problems benefit from intrinsic simplification

III. Intrinsic simplification

- Extrinsic methods preserve irrelevant extrinsic details
- Intrinsic approach opens up a larger space of triangulations
- Extreme example: neardevelopable surfaces

intrinsic simplification


## Inspiration: quadric error simplification

[Garland \&
Heckbert 1997]
III. Intrinsic simplification - motivation

1. Local simplification operation

2. Accumulated distortion measurements


- Algorithm: repeatedly collapse cheapest edge
- Efficient: all local operations
- Accurate: accumulates error estimates


## Intrinsic simplification

1. Local simplification operation

intrinsic vertex removal
2. Accumulated distortion measurements

intrinsic curvature error

- Algorithm: repeatedly remove cheapest vertex


## Intrinsic simplification

III. Intrinsic simplification

1. Local simplification operation

intrinsic vertex removal
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- Algorithm: repeatedly remove cheapest vertex


## Intrinsic vertex removal

- Intrinsic view: replace curved vertex with flat patch
- i.e., parameterization problem



## Intrinsic vertex removal

- Intrinsic view: replace curved vertex with flat patch
- i.e., parameterization problem



## Vertex flattening

- Map 1-ring to plane such that:
(1) Distortion is low
(2) Boundary edge lengths are preserved
III. Intrinsic simplification - intrinsic vertex removal

- Discrete conformal map [Springborn, Schröder \& Pinkall 2008]
- Fix $u=0$ on boundary
- Efficient 1D optimization problem


## Flat vertex removal

- Same as before



## Intrinsic simplification

III. Intrinsic simplification - intrinsic curvature error

1. Local simplification operation

intrinsic vertex removal
2. Accumulated distortion measurements

intrinsic curvature error

- Algorithm: repeatedly remove cheapest vertex


## Distortion: curvature redistribution



We approximate the transport cost of this curvature redistribution

## Mass transport cost



Vertex removal


## Comparing mass distributions


III. Intrinsic simplification - intrinsic curvature error

Goal: approximate cost of transporting each distribution to its vertex
mass distribution transported to vertex $i$
mass distribution transported to vertex $j$

## Approximating the mass transport cost

III. Intrinsic simplification


## Specializing to curvature

- Challenge: curvature is signed
- Just track positive and negative parts separately


## Simplification with the curvature transport cost

III. Intrinsic simplification - intrinsic curvature error


## Other transport costs

III. Intrinsic simplification - intrinsic curvature error

- Track transport of other data (e.g. area) in same way
- Can take weighted combinations of costs



## Surface correspondence

- Simplified mesh not isometric to original surface

- But, only uses a few local operations
- Correspondence easy for each operation
- Encode correspondence via list of operations

[^1]

## Prolongation

- Transfer piecewise-linear functions:
- Just find values at vertices
- Encode by a matrix



## Pulling back vector fields

- correspondence
- Approximate differential of correspondence map

Encode by complex prolongation matrix


## Intrinsic simplification-summary

III. Intrinsic simplification

- summary

1. Local simplification operation

intrinsic vertex removal
2. Accumulated distortion measurements

intrinsic curvature error

- Algorithm: repeatedly remove cheapest vertex
- Correspondence: record operation history


## Results

## Surface hierarchies



## Hierarchies accelerate computation

- Accelerate many geometric tasks
- Even helps with extrinsic problems



## Distortion


i.e. quasiconformal dilatation
(mean 1.115)
(mean 8.1\%)

## Geodesic distance

III. Intrinsic simplification


## Geodesic Voronoi diagrams

III. Intrinsic simplification
result on
simplified surface
ground truth
7207.4 ms
only $1 \%$ vertices misclassified

## Speedup vs error in geodesic distance


speedup/error: 3x/0.0002\%


840x / 0.2\%


## Low rank all-pairs distance matrix approximation


III. Intrinsic simplification


Distance matrix of simplified mesh
$\overbrace{}^{P}$
Prolongation operator

- Approximate distance matrix
$P: \mathbb{R}^{\left|V^{c}\right|} \rightarrow \mathbb{R}^{|V|}$
$\hat{D}=P \tilde{D} P^{\top}$


## Adaptive simplification


III. Intrinsic simplification
results


## input


input
constrained coarsening
Poisson solve


## Performance

- Linear scaling
- Constant work per vertex

Removes ~10,000 vertices per second
time (s)
II. Intrinsic simplification


## Try it out yourself (... in the near future)



## IV. Discrete Uniformization



[
G., Springborn, \& Crane. 2021. Discrete conformal equivalence of polyhedral surfaces. ACM Transactions on Graphics

## The uniformization theorem

 [Poincare 1907; Koebe 1907; Troyanov 1991]
IV. Discrete uniformization

Any surface is conformally equivalent to a surface of constant curvature.


Image: [Crane, Pinkall \& Schröder 2013]

## The discrete uniformization theorem

[Gu, Luo, Sun \& Wu 2018; Springborn 2019]

IV. Discrete uniformization

Any positive vertex cone angles satisfying GaussBonnet can be realized by some discrete conformal map.


## The discrete spherical uniformization theorem

[Springborn 2019]
IV. Discrete uniformization

Any simply-connected triangle mesh is discretely conformally equivalent to a mesh whose vertices lie on the unit sphere


# Discrete uniformization in action 

[G., Springborn, \& Crane. 2021]


## Triangle mesh $\hookleftarrow$ ideal polyhedron

[Bobenko, Pinkall \& Springborn 2010]

IV. Discrete uniformization


Euclidean triangle in circumcircle $\hookleftarrow$ Klein ideal triangle

## Ideal Delaunay triangulations



Hyperbolic correspondence problem

## Correspondence between ideal polyhedra

- Adapt Euclidean techniques to hyperbolic setting
- Integer coordinates essential



## Projective interpolation

- [Springborn, Schröder \& Pinkall 2008]: projective interpolation
- Hyperbolic isometry
- In variable triangulation case, lay out triangles in hyperboloid model



## Interpolation in the hyperboloid model

IV. Discrete uniformization


## Interpolation in the hyperboloid model


fixed triangulation

variable triangulation

## Starting from Delaunay



## Final algorithm



flip to (Euclidean)
Delaunay

find scale factors

lay out in plane
$\qquad$
extract correspondence

compute common subdivision
$\qquad$
interpolate via hyperboloid

## Uniformization results

## Challenging datasets

IV. Discrete uniformization

- results
difficult cones


Dataset

## MPZ

| [Myles+ 2014] | 114 | $100 \%$ | 8 s |
| :---: | :---: | :---: | :---: |
| Thingi10k <br> [Zhou+2016] | $32,744^{*}$ | $97.7 \%$ | $57 \mathrm{~s}^{\dagger}$ |
| brain scans <br> $[Y e o+2009]$ | 78 | $100 \%$ | 493 s |
| anatomical surfaces <br> [Boyer+ 2011] | 187 | $100 \%$ | 15 s |

## Variable triangulation > fixed triangulation

IV. Discrete uniformization


Fixed triangulation (CETM)


## Boundary conditions


IV. Discrete uniformization - results
convex


## Try it out yourself


IV. Discrete uniformization

- results


C++ application
Save textured mesh
projective interpolation in Blender

https://github.com/MarkGillespie/CEPS

## Thanks for listening



- And thank you to all of my great coauthors!


Boris
Springborn


Keenan
Crane


Nicholas
Sharp


Hsueh-Ti
Derek Liu


Benjamin Chislett


Alec Jacobson

## Supplemental Slides

## The importance of memory



Memoryless transport cost
Vertex removal


## Robust hierarchies



## Tangent space approximations



## Signed curvature transport



$$
\alpha_{i j}:=\frac{\left|\tilde{K}_{j}-K_{j}\right|}{\sum_{l \in \mathcal{N}_{i}}\left|\tilde{K}_{l}-K_{l}\right|}
$$


[^0]:    ( concretely, just 3 integers per mesh edge )

[^1]:    1. Flip edge 1
    2. Scale vertex 5
    3. Remove vertex 5
    4. Flip edge 8
    5. Flip edge 12
    6. Scale vertex 2
